# NEW SOLUTIONS OF $a^{p-1} \equiv 1\left(\bmod p^{2}\right)$ 

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Dedicated to the memory of my undergraduate advisor, D. H. Lehmer

Abstract. We tabulate solutions of $a^{p-1} \equiv 1\left(\bmod p^{2}\right)$ where $2 \leq a \leq 99$ and where $p$ is an odd prime, $p<2^{32}$.

## 1. Introduction and summary

Some number-theoretic questions such as Fermat's conjecture [4] require primes $p$ satisfying

$$
\begin{equation*}
a^{p-1} \equiv 1 \quad\left(\bmod p^{2}\right) \tag{1}
\end{equation*}
$$

for given $a$ not a power. Brillhart, Tonascia, and Weinberger [2] list all solutions of (1) for $2 \leq a \leq 99$ and $3 \leq p<10^{6}$, plus some solutions for higher $p$. Lehmer [3] subsequently extended the $a=2$ search to $p<6 \cdot 10^{9}$, finding only the known solutions $p=1093$ and $p=3511$. Aaltonen and Inkeri [1] list solutions for prime $a<1000$ and $p<10^{4}$. Table 1 (next page) extends the table in [2] to $p<2^{32}$, giving 23 new solutions. Included are the first solutions for $a=66$ and $a=88$.

The table in [2] identifies where (1) holds modulo $p^{3}$, with the only solutions for $a \leq 99$ and $p>7$ being $(a, p)=(42,23)$ and $(68,113)$. This search found no more such solutions.

The pair $(a, p)=(5,1645333507)$ satisfies $p^{a-1} \equiv 1\left(\bmod a^{2}\right)$ as well as (1). This supplements the pairs $(2,1093),(3,1006003)$, and $(83,4871)$ listed in [1, p. 365].

The largest known $p$ for which multiple $a$ satisfy (1) with $2 \leq a \leq 99, a$ not a power, is $p=331$, for which $a=18$ and $a=71$ satisfy (1).

The Fibonacci sequence is defined by $F_{0}=0, F_{1}=1$, and $F_{n}=F_{n-1}+F_{n-2}$ for $n \geq 2$. If $p \neq 5$, then $F_{p-\epsilon} \equiv 0(\bmod p)$, where $\epsilon=+1$ if $p \equiv \pm 1$ $(\bmod 5)$ and $\epsilon=-1$ if $p \equiv \pm 2(\bmod 5)$. Williams [5, pp. 85-86] reports no solution of $F_{p-\epsilon} \equiv 0\left(\bmod p^{2}\right)$ with $p<10^{9}$. This search found no such solution with $p<2^{32}$.

[^0]Table 1. Solutions of $a^{p-1} \equiv 1\left(\bmod p^{2}\right)$ with $2 \leq a \leq 99$ and $3 \leq p<2^{32}$. New solutions are in bold font

| $a$ | Values of $p$ | $a$ | Values of $p$ |
| :---: | :---: | :---: | :---: |
| 2 | 10933511 | 55 | 3301097278001 |
| 3 | 111006003 | 56 | 6477079771 |
| 5 | 2077140487534711611645333507 | 57 | 54769986197 |
| 6 | 661615348513152573 | 58 | 13142250279 |
| 7 | 5491531 | 59 | 2777 |
| 10 | 348756598313 | 60 | 29 |
| 11 | 71 | 61 |  |
| 12 | 2693123653 | 62 | $\begin{array}{lllll}3 & 19 & 127 & 1291\end{array}$ |
| 13 | 8631747591 | 63 | $\begin{array}{lllllll}23 & 29 & 36713 & 401771\end{array}$ |
| 14 | 29353 | 65 | 17163 |
| 15 | 29131 | 66 | 89351671 |
| 17 | 34602148947 | 67 | 747268573 |
| 18 | $\begin{array}{llllllll}5 & 7 & 37 & 331 & 339231284043\end{array}$ | 68 |  |
| 19 | $\begin{array}{llllllll}3 & 7 & 13 & 43 & 137 & 63061489\end{array}$ | 69 | $\begin{array}{llllll}19 & 223 & 631 & 2503037\end{array}$ |
| 20 | 281464579377747122959073 | 70 | 13142963 |
| 21 |  | 71 | 347331 |
| 22 | $13 \quad 6731595813492366587$ | 72 |  |
| 23 | $13 \quad 248175713703077^{1}$ | 73 | 3 |
| 24 | 525633 | 74 | 5 |
| 26 | $3 \mathrm{~S}^{5} 71486999673$ | 75 | $\begin{array}{lllll}17 & 43 & 347 & 31247\end{array}$ |
| 28 | $319 \quad 23$ | 76 | 515711099241661049 |
| 29 |  | 77 | 32687 |
| 30 | 7160541 | 78 | $\begin{array}{lllllllll}43 & 151 & 181 & 1163 & 56149 & 4229335793\end{array}$ |
| 31 | 77964512806861 | 79 | $\begin{array}{llllll}7 & 263 & 3037 & 1012573 & 60312841\end{array}$ |
| 33 | 23347441 | 80 | $\begin{array}{lllll}3 & 7 & 13 & 6343\end{array}$ |
| 34 |  | 82 | 35 |
| 35 | 316133571 | 83 | 487113691315746063 |
| 37 | 377867 | 84 | 16365320101 |
| 38 | 17127 | 85 | 11779 |
| 39 | 8039 | 86 | 68239 |
| 40 | $\begin{array}{llllll}11 & 17 & 307 & 66431\end{array}$ | 87 | 199948121 |
| 41 | 291025273138200401 | 88 | 2535619637 |
| 42 | 23 | 89 | 313 |
| 43 | 5103 | 90 |  |
| 44 | 32295851 | 91 | 3293 |
| 45 | 1283131759157635607 | 92 | $\begin{array}{llllllll}727 & 383951 & 12026117 & 18768727 & 1485161969\end{array}$ |
| 46 | 3829 | 93 | $5 \quad 509922181551$ |
| 47 |  | 94 | 11124132143463033 |
| 48 | 7257 | 95 | 213715061 |
| 50 | 7 | 96 | 1095437832912925267 |
| 51 | 541 | 97 | 72914393 |
| 52 | 4611228488439 | 98 | 32862761001527 |
| 53 | $3 \quad 47 \quad 5997$ | 99 | $\begin{array}{llllll}5 & 7 & 13 & 19 & 83\end{array}$ |
| 54 | 191949 |  |  |

${ }^{1}$ Incorrectly printed as "1370377" in [2].

## 2. Programming considerations

As in [2] and [3], it suffices to compute the last two digits of the base $p$ representation of each intermediate result. Since (1) is equivalent to $a^{(p-1) / 2} \equiv$ $\pm 1\left(\bmod p^{2}\right)$, we can save a squaring $\bmod p^{2}$.

The programs in [2] fixed the base $a$ and looped through values of $p$. One can instead check all values of $a$ together for a given $p$. Then the value of $a^{(p-1) / 2}\left(\bmod p^{2}\right)$ need be calculated the long way (binary method of exponentiation) only for prime $a$ : if $a=a_{1} a_{2}$ where

$$
a_{1}^{(p-1) / 2} \equiv \pm\left(1+p b_{1}\right) \quad\left(\bmod p^{2}\right) \quad \text { and } \quad a_{2}^{(p-1) / 2} \equiv \pm\left(1+p b_{2}\right) \quad\left(\bmod p^{2}\right)
$$

then $a^{(p-1) / 2} \equiv \pm\left(1+p\left(b_{1}+b_{2}\right)\right)\left(\bmod p^{2}\right)$. The latter computation reduces to an addition modulo $p$. Since $\pi(100)=25$ whereas there are 87 nonpowers below 100 , this represents a potential $70 \%$ savings.

The search for $p<2^{31}$ was done on a DECstation 3100 (MIPS architecture). To compute a product $a b \bmod p$ where $0 \leq a, b<p$ but where $a b$ may exceed the largest single-precision integer, the program computed $q=a \cdot b \cdot \frac{1+\epsilon}{p}$, using floating-point arithmetic, where $2^{-50} \ll \epsilon \ll 1 / p$. The relative error in any floating-point computation is at most $2^{-52}$ (53-bit mantissas), ensuring that

$$
\frac{a b}{p} \leq q \leq \frac{a b}{p}(1+1 / p)<\frac{a b}{p}+1
$$

and hence that $\left\lfloor\frac{a b}{p}\right\rfloor \in\{\lfloor q\rfloor,\lfloor q\rfloor-1\}$; the choice is made using the sign of $r=a b-p\lfloor q\rfloor$. Since $-2^{31}<-p \leq r<p<2^{31}$, this $r$ can be computed by integer arithmetic modulo $2^{32}$.

This technique fails for $p>2^{31}$ unless the program uses 64-bit arithmetic to compute the tentative remainder (it would also require converting unsigned 32bit integers to/from floating point). Instead, the computations for $p>2^{31}$ were done on a NeXT with a Motorola 68040 chip. The 68040 can divide a 64 -bit unsigned integer by a 32-bit unsigned integer, obtaining quotient and remainder in one instruction (if the quotient does not overflow), but the MIPS architecture lacks such. The 3100 tried all primes in an interval of length 10 million per hour. The 68040 computations took slightly longer, searching an interval of length 7 million per hour.

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